

Research and Development of Digital Signatures

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Abstract: Since Diffie and Hellman's pioneering work on asymmetric cryptography in 1976, digital signature technology has evolved through three phases—theoretical foundation, standardization, and diversified innovation—emerging as a cornerstone of trust in digital societies. Theoretically, foundational frameworks were established by RSA, DSA, and Schnorr algorithms. Standardization efforts, including NIST DSS, ISO/IEC series, and national systems (e.g., China's SM2/SM9, Russia's GOST), fostered a multipolar ecosystem. Extended-attribution technologies (blind, group, and ring signatures) addressed privacy and scenario-specific demands. Current challenges, such as quantum computing threats and privacy-regulation trade-offs, drive advancements in post-quantum cryptography (lattice-based signatures, hash-based XMSS) and privacy-enhancing mechanisms (verifiably encrypted signatures, homomorphic signatures), guided by ISO/IEC redactable standards and NIST's post-quantum initiative. Moving forward, digital signatures will deepen capabilities in provable security, quantum resistance, and adaptive policy control, underpinning trust architectures for emerging ecosystems like Web3 and the metaverse.

Keywords: Digital Signature Technology; Privacy-Enhancing Mechanisms; Standardization Process.

1. Introduction

Traditional authentication methods such as handwritten signatures, seals, and biometric features (e.g., fingerprints) have historically served as fundamental mechanisms for identity verification and document authentication in human societal development. These physical validation techniques operate by establishing traceable individual identifiers, thereby enabling three critical security properties: authenticity (confirming the identity of the document signer), integrity (ensuring the unaltered state of document content), and non-repudiation (preventing the signer from denying their participation in the signing act). With the advancement of digital transformation, digital signature technology has emerged as a seminal achievement in cryptography. Leveraging asymmetric cryptographic mechanisms, it not only inherits the core functionalities of physical signatures but also achieves substantial security augmentation. This innovation addresses inherent vulnerabilities in traditional methods, such as susceptibility to forgery and limited traceability, while maintaining compliance with the tripartite security requirements of authenticity, integrity, and non-repudiation in electronic environments.

The core architecture of digital signature technology is based on the Public Key Infrastructure (PKI), with its technical characteristics manifested in three dimensions: firstly, ensuring the authenticity of the signer's identity through the unique binding mechanism between digital certificates and private keys; secondly, guaranteeing data integrity by generating message digests using cryptographic hash functions combined with digital signature algorithms; thirdly, achieving non-repudiation of signing acts through mathematically irreversible cryptographic principles. Compared to traditional signature methods, digital signatures demonstrate significant technical advantages in implementation: the signature generation phase requires encryption operations using the signer's exclusive private key, while the verification process depends solely on publicly distributed public keys. This asymmetric mechanism effectively addresses technical flaws such as susceptibility to forgery and difficulty in traceability inherent in traditional

signatures.

From a technical implementation perspective, a standard digital signature system comprises the following core components: (1) a key generation algorithm that produces mathematically linked public-private key pairs; (2) a signing algorithm that cryptographically binds the private key to a message digest; (3) a verification algorithm that decrypts and validates the signature using the public key. For large-scale data scenarios, the system adopts a preprocessing mechanism—generating fixed-length message digests via collision-resistant hash functions before performing signature operations on the digests—thereby ensuring both computational efficiency and cryptographic security requirements are met while maintaining compliance with the fundamental principles of digital signature protocols.

The security of digital signature mechanisms must satisfy the following fundamental requirements: first, the computational infeasibility of deriving the private key from the public key within polynomial time (i.e., the one-way property of the trapdoor function); second, the inability of unauthorized entities to construct valid message-signature pairs. According to modern cryptographic security models, the mechanism must achieve Existential Unforgeability under Adaptive Chosen-Message Attacks (EUF-CMA) [1], meaning that even after obtaining polynomially bounded message-signature pairs, an adversary cannot forge a valid signature for a new message. The more stringent Strong Existential Unforgeability (SUF-CMA) [2] further requires that an adversary cannot generate a new valid signature for an already signed message, a property critical for defending against replay attacks and signature splitting attacks.

2. Evolution of Digital Signatures

Since Whitfield Diffie and Martin Hellman introduced the concept of digital signatures in their seminal 1976 paper *New Directions in Cryptography* [3], this technology has undergone over four decades of theoretical breakthroughs and practical evolution. From early feasibility explorations rooted in computational complexity theory, to the formalization of security models and the realization of provably secure

constructions, and further to the development of feature-specific technological advancements for diverse application scenarios, digital signature technology has matured from an academic concept into a core security infrastructure underpinning modern digital societies.

Key milestones in its developmental trajectory include: (1) Diffie and Hellman's pioneering proposal of the asymmetric cryptographic framework in 1976, establishing the theoretical foundation for digital signatures [3]; (2) Rivest, Shamir, and Adleman's introduction of the first practical RSA algorithm in 1978 [4], accompanied by Rabin's modular square-based one-time signature scheme the same year [5]; (3) Merkle's 1979 hash chain-based signature algorithm [6], whose Merkle tree structure became a foundational component for blockchain technologies; (4) Elgamal's discrete logarithm-based signature framework [9], Shamir's identity-based cryptography paradigm (IBC) [10], and Goldwasser et al.'s formalization of the EUF-CMA security model [1] in 1984; (5) The Fiat-Shamir transformation enabling systematic conversion of authentication protocols to digital signatures in 1986 [11]; (6) Schnorr's modular signature scheme [12] and the first commercial implementation in Lotus Notes 1.0 [14], both emerging in 1989.

The standardization process accelerated in the 1990s with critical advancements: (7) the 1991 NIST publication of the DSA standard [15] and Zimmermann's release of PGP 1.0 supporting RSA signatures [17]; (8) the 1993 PKCS#1 v1.5 specification standardizing RSA encoding rules [20]; (9) the 1995 integration of signature algorithms into the SSL 2.0 protocol [21]; (10) Bellare and Rogaway's 1996 provably secure RSA-PSS scheme [22], Pointcheval and Stern's formalization of the Forking Lemma for security analysis [23], and Kocher's groundbreaking revelation of side-channel attack vulnerabilities [24]; (11) Gennaro et al.'s 1999 realization of RSA security constructions in the Standard Model [25, 26], marking a paradigm shift in cryptographic proof methodologies.

Since the 21st century, digital signature technology has diversified across multiple directions: (12) The 2001 introduction of the BLS short signature algorithm [27], which significantly improved spatial efficiency through pairing-based cryptography; (13) The successive proposals of identity-based signature (IBS) schemes leveraging bilinear pairings in 2002 [28, 29]; (14) The 2003 Al-Riyami-Paterson certificateless signature framework [30], initiating research into novel paradigms beyond traditional PKI; (15) The 2008 breakthrough by Gentry et al. in achieving provably secure lattice-based signatures [34], laying groundwork for post-quantum cryptography; (16) The 2017 launch of NIST's post-quantum cryptography standardization initiative [37], accelerating the development of quantum-resistant algorithms like XMSS; (17) The 2018 unification of certificateless signature security models by Cheng and Chen [31], resolving critical inconsistencies in prior frameworks.

This evolutionary trajectory reveals two fundamental principles: cryptographic theoretical breakthroughs (e.g., bilinear pairings, lattice theory) persistently catalyze novel signature paradigms, while evolving security threats (e.g., quantum computing, side-channel attacks) continuously drive technological innovation. Currently, digital signature technology exists in a transitional phase of coexistence between classical security models and post-quantum cryptographic frameworks. As quantum-resistant algorithms mature and hybrid systems emerge, its ongoing evolution will

persistently reshape the trust architecture of digital societies, balancing legacy compatibility with forward-looking cryptographic resilience.

3. Multifunctional Digital Signatures

With the maturation of generic digital signature technologies, numerous extended-attribution digital signature schemes have emerged to address privacy preservation and fairness requirements in specialized application scenarios such as electronic cash and electronic voting. The following enumerates representative technical paradigms and their core characteristics (this list serves as an illustrative subset; comprehensive analyses may be found in systematic surveys such as [38, 39]):

1) Blind Signatures: Proposed by Chaum in 1982 [40], this mechanism permits signers to generate valid signatures on blinded messages without accessing plaintext content, enabling anonymous authentication critical for electronic voting and digital currency systems.

2) Multi-Signatures & Aggregate Signatures: Introduced by Itakura-Nakamura in 1983 [41], multi-signatures allow multiple signers to collaboratively produce a compact signature, while aggregate signatures [42, 43] enable compression of multiple independent signatures into a single verifiable unit, optimizing blockchain transaction validation.

3) Threshold Signatures: Formalized by Desmedt in 1987 [44], a (k,n) -threshold scheme mandates that any subset of k signers from n participants must collaborate to generate valid signatures, enhancing robustness against insider threats.

4) Undeniable Signatures: Designed by Chaum-van Antwerpen in 1989 [45], this interactive mechanism requires signer participation during verification, preventing third-party misuse of signed data.

5) Fail-Stop Signatures: Conceptualized by Pfitzmann in 1991 [46], this paradigm provides cryptographic evidence to expose forgery attempts, offering enhanced protection against quantum adversaries.

6) Group Signatures: Proposed by Chaum-van Heyst in 1991 [47], this framework enables group members to sign anonymously while permitting authorized tracing via a group manager, balancing privacy and accountability in consortium blockchains.

7) Designated Confirmer Signatures: Introduced by Chaum in 1994 [50], this semi-trusted third-party mechanism delegates signature verification authority to a designated confirmer via verifiable encryption. The scheme preserves signer control while resolving verification deadlocks caused by signer non-cooperation in undeniable signature protocols.

8) Proxy Signatures: Proposed by Mambo et al. in 1996 [51], this delegation model enables secure transfer of signing rights through digital authorization credentials. Proxy signers gain limited signing privileges without accessing the original private key, making it suitable for permission management in distributed systems.

9) Designated Verifier Signatures: Developed by Jakobsson et al. in 1996 [52], this framework binds verification capability to a specific recipient using cryptographic key pairing, ensuring only designated parties can validate signatures. By eliminating third-party verification, it enhances privacy in scenarios like confidential business contracts.

10) Signcryption: Pioneered by Yuliang Zheng in 1997 [53], this cryptographic hybrid integrates encryption and signing

into a single operation, reducing computational overhead by 50-70% compared to traditional "sign-then-encrypt" approaches. Its variants now include identity-based [80] and post-quantum secure [81] implementations, particularly advantageous for IoT and 5G networks.

11) Ring Signatures & Derivatives: Rivest et al.'s 2001 framework [54] enables ad-hoc anonymous signing from arbitrary public key sets (rings) without group administrators. Key variants include: Threshold ring signatures [55] enforcing (k,n) collaborative signing policies; Linkable ring signatures [56] detecting duplicate signer activity while preserving anonymity; Traceable ring signatures [57] balancing auditability through controlled identity revelation; Mesh signatures [58] enabling non-PKI-based group expansion, widely adopted in privacy coins like Monero.

12) Signatures with Efficient Protocols: Camenisch and Lysyanskaya's 2001 paradigm [59] combines secure multi-party computation with zero-knowledge proofs to enable: Blind signature generation with content privacy preservation; Efficient proof-of-signature-possession mechanisms; This breakthrough underpins critical applications like Direct Anonymous Attestation (DAA) [60] in trusted platform modules.

13) Homomorphic Signatures: Johnson et al.'s 2002 innovation [61] allows computations (linear combinations [63], polynomial transforms [64]) on signed data while preserving verifiability. Extended implementations include: Sanitizable signatures [67] for authorized partial modifications; Append-only signatures [68] supporting dynamic data expansion; Blank signatures [69] enabling template-based signing workflows; Protean signatures [70] integrating multi-dimensional revisions; Standardization efforts under ISO/IEC 29167-20 [71] aim to unstrate these for industrial adoption.

14) Verifiably Encrypted Signatures: Boneh et al.'s 2003 construction [72] employs arbitrator-controlled encryption to ensure fair contract execution. Its evolution, commuting signatures, enables encrypted-domain verification, overcoming limitations of separate encryption-signature architectures.

15) Concurrent Signatures: Chen et al.'s 2004 protocol [73] introduces keystone-controlled atomic exchange, ensuring mutual signature binding in e-commerce transactions while preventing post-signing repudiation.

16) Anonymous Signatures: Yang et al.'s 2006 design [74] decouples signer identity from message context, achieving metadata-free verification for anonymous credential systems.

17) Signatures of Knowledge: Chase and Lysyanskaya's 2006 framework [75] embeds zero-knowledge proofs into signatures, enabling implicit assertions (e.g., group membership) without revealing evidence.

18) Structure-Preserving Signatures: Abe et al.'s 2010 scheme [76] constrains cryptographic elements to bilinear groups, enabling modular composition with privacy-preserving protocols like anonymous credentials.

19) Attribute-Based Signatures: Maji et al.'s 2011 model [77] binds signing rights to attribute sets, with policy-based extensions [78] enabling complex logical expressions for cloud access control.

20) Functional Signatures: Boyle et al.'s 2014 framework [79] restricts signing to function outputs $f(m)$, establishing foundational cryptography for verifiable cloud computing and federated learning.

These extended signature technologies share a common

characteristic: the innovative integration of cryptographic primitives. This integration enables functional extensions tailored to specific requirements—such as privacy enhancement, access control, and efficiency optimization—while preserving core authentication capabilities. This evolutionary trajectory exemplifies the "demand-driven innovation" principle in cryptography, demonstrating how theoretical advancements respond to real-world challenges. By systematically addressing diverse needs across digital ecosystems, these technologies collectively furnish a versatile toolkit for constructing sophisticated digital trust architectures, balancing foundational security guarantees with adaptive functionality for emerging use cases.

4. Standardization Journey

The standardization process of digital signature technologies has evolved into a multipolar landscape with diversified developments. In the United States, NIST published the Digital Signature Algorithm (DSA) in 1991 and formally established it as the Digital Signature Standard (DSS) through FIPS 186 in 1994 [15]. The elliptic curve variant ECDSA achieved standardization via FIPS 186-2 in 2000, forming a dual-track technical framework combining discrete logarithm and elliptic curve cryptography. The standardization of RSA reflects an evolutionary security philosophy: PKCS#1 v1.5 (1993) adopted the EMSA-PKCS1-v1_5 padding scheme [20], while v2.2 introduced the provably secure RSA-PSS scheme in 2002, later solidified at the protocol layer through RFC 8017 in 2016 [81]. In 2017, the EdDSA algorithm, leveraging Edwards-curve advantages and formal security proofs, was standardized in RFC 8032 [82], ultimately contributing to TLS 1.3 protocol [83] establishing its quadruple algorithm system: RSASSA-PKCS1-v1_5, RSA-PSS, ECDSA, and EDDSA.

The international standardization landscape exhibits a diversified competitive pattern: Russia established an independent cryptographic system through the GOST R 34.10 series standards (1994-2012) [84-86]; regional standards such as South Korea's KCDSA/ECKCDA (1998) [87-89], Germany's EC-GDSA (2005) [90], and China's SM2/SM9 (2012-2016) [91-92] were subsequently formalized. ISO/IEC has constructed a multidimensional standardization framework encompassing five technical dimensions: Message Recovery Mechanisms: ISO/IEC 9796 series standards [93-94] governing digital signatures with message recovery capabilities. Appendix-Type Signatures: ISO/IEC 14888 series standards [95-96], where Part 2 specifies seven integer factorization-based algorithms and Part 3 integrates fourteen discrete logarithm-based schemes. Anonymous Signatures: ISO/IEC 20008 series standards [97] defining group public key-based anonymous signing mechanisms. Blind Signatures: ISO/IEC 18370 series standards [98] standardizing discrete logarithm-based blind signature implementations. Redactable Signatures: ISO/IEC 23264 series standards [71] regulating asymmetric redaction mechanisms for authenticated data.

The ISO/IEC 13888 non-repudiation technical standard system [99-100] establishes a comprehensive assurance framework encompassing source authentication, delivery confirmation, and timestamp validation through dual-track symmetric and asymmetric cryptographic mechanisms. Its technical implementation pathways include a composite verification mechanism based on symmetric cryptography (ISO/IEC 13888-2) and multi-modal non-repudiation

protocols leveraging digital signatures (ISO/IEC 13888-3), providing systematic protection against repudiation risks across transactional lifecycle phases.

5. Summary

Digital signature technology has evolved over four decades into a mature ecosystem that integrates theoretical innovation, standardization, and industrial application. This article systematically reviews key technological breakthroughs, characteristic evolutionary trajectories, and standardization processes, revealing its transformation from cryptographic theoretical achievements to a trust cornerstone of digital society. With emerging threats such as quantum computing, digital signature technology will continue advancing while preserving foundational security attributes, deepening developments in provable security, quantum resistance, privacy enhancement, and related directions.

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